



PRESIDENCY UNIVERSITY BENGALURU

Roll No. _____

SCHOOL OF ENGINEERING

TEST - 1

Sem: Odd Sem 2019-20

Date: 01-09-2019

Course Code: CSE 208

Time: 11AM to 12PM

Course Name: THEORY OF COMPUTATION

Max Marks: 40

Program & Sem: B.Tech, (CSE) & III

Weightage: 20%

Instructions:

- (i) Answer all the questions from each part

Part A [Memory Recall Questions]

Answer all the Questions. Each Question carries two marks. (5Qx2M=10M)

1. Define the Language of an automata with an example. (C.O.NO.1)[Knowledge]
 2. Define the following with an example for each.
 - a. Strings
 - b. Alphabet(C.O.NO.1)[Knowledge]
 3. Describe the deterministic finite automata with five component values. (C.O.NO.1)[Knowledge]
 4. List the applications of finite automata. (C.O.NO.1)[Knowledge]
 5. Describe the transition diagram with an example. (C.O.NO.1)[Knowledge]

Part B [Thought Provoking Questions]

Answer both the Questions. Each Question carries five marks. (2Qx5M=10M)

6. Shreyas is working on an abstract machine design. He wants to design a deterministic machine in such a way that only the strings with even occurrence of 1's and odd occurrence of 0's must be accepted by the accepter. Construct the DFA for the above scenario and provide the transition diagram, transition table, five component values. Using extended transition function show the acceptance of the string "10010"

(C.O.NO.2.)[Comprehension]

7. Anu is having a collection of symbols with a's and b's. A deterministic automata has to be designed to accept the strings having "aa" but not "aaa". The occurrence of the symbol 'b' can be anywhere in the path of the string. The solution to the above problem must be provided with transition diagram, transition table, five component values and also show one valid string acceptance using extended transition function.

(C.O.NO.2.)[Comprehension]

Part C [Problem Solving Questions]

Answer both the Questions. Each Question carries ten marks. (2Qx10M=20M)

8. Convert the following Non Deterministic finite automata into deterministic finite automata. (C.O.NO.2)[Application]

	0	1
$\rightarrow p$	{p, q}	{p}
q	{r}	{r}
r	{s}	\emptyset
*s	{s}	{s}

9. Convert the following ϵ -NFA to its equivalent DFA. (C.O.NO.2)[Application]





SCHOOL OF ENGINEERING

Semester: 3rd

Course Code: CSE 208

Course Name: Theory of Computation

Date: 01/10/2019

Time: 11AM to 12PM

Max Marks: 40

Weightage: 20%

Extract of question distribution [outcome wise & level wise]

Q.NO	C.O.NO	Unit/Module Number/Unit /Module Title	Memory recall type [Marks allotted]		Thought provoking type [Marks allotted]		Problem Solving type [Marks allotted]		Total Marks
			Bloom's Levels	K	Bloom's Levels	C	Bloom's Levels	A	
1	1	Module -1		02					02
2	1	Module -1		02					02
3	1	Module -2		02					02
4	1	Module - 1		02					02
5	1	Module - 2		02					02
6	2	Module -2				05			05
7	2	Module -2				05			05
8	2	Module -2						10	10
9	2	Module -2						10	10
	Total Marks			10		10		20	40M

I hereby certify that all the questions are set as per the guide lines [Manipal B.C].

Reviewer's Comments:

R

Annexure- II: Format of Answer Scheme



SCHOOL OF ENGINEERING

SOLUTION

Date: 01/10/2019

Time: 11AM to 12PM

Max Marks: 40

Weightage: 20%

Semester: 3rd

Course Code: CSE 208

Course Name: Theory of Computation

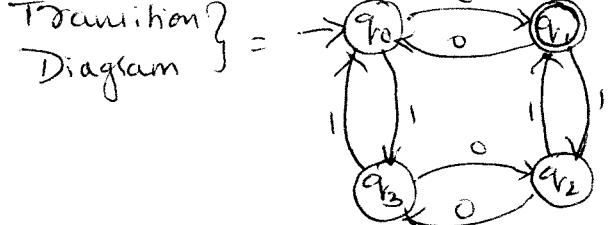
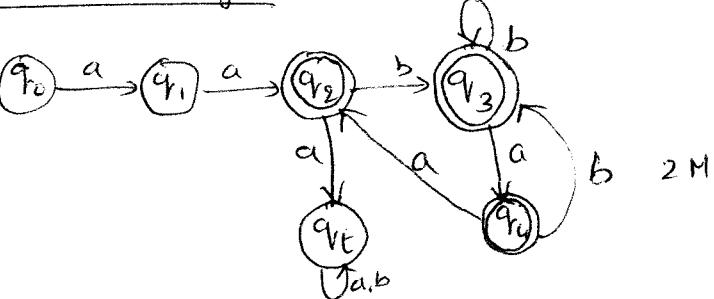
Part A

(5 x 2 =10 Marks)

Q No	Solution	Scheme of Marking	Max. Time required for each Question
1.	The set of all strings from some Σ^* is called the language of an automata. Eq: $L = \{w : w \in \{a,b\}^*\}$	2 M	2 min
2.a	An alphabet is a finite set of symbols. It is denoted as Σ . Eq: $\Sigma = \{a, b\}$, $\Sigma = \{0, 1\}$	2 M (1+1)	2 min
b.	Finite sequence of symbols taken from some Σ is called the string. If $\Sigma = \{0, 1\}$, The string = $\{0, 1, 001, 010, \dots\}$		
3.	DFA $M = \{Q, \Sigma, \delta, q_0, F\}$ Define each tuple	2 M	2 min
4.	List atleast 4 applications of finite automata	2 M	2 min
5.	With an example describe about Transition diagram	2 M	2 min

Part B

(2x 5 = 10 Marks)

Q No	Solution	Scheme of Marking	Max. Time required for each Question																
6.	<p>The alphabet recognized based on the given scenario is $\Sigma = \{0, 1\}$.</p> <p>Transition Diagram = </p> <p>Transition Table</p> <table border="1" data-bbox="589 675 886 967"> <thead> <tr> <th>States</th> <th>0</th> <th>1</th> </tr> </thead> <tbody> <tr> <td>q_0</td> <td>q_1</td> <td>q_3</td> </tr> <tr> <td>q_1</td> <td>q_0</td> <td>q_2</td> </tr> <tr> <td>q_2</td> <td>q_3</td> <td>q_0</td> </tr> <tr> <td>q_3</td> <td>q_2</td> <td>q_0</td> </tr> </tbody> </table> <p>Five components $\{\emptyset, \Sigma, \delta, q_0, F\}$ values</p> <p>'10010' string derivation using δ^*</p>	States	0	1	q_0	q_1	q_3	q_1	q_0	q_2	q_2	q_3	q_0	q_3	q_2	q_0	2 M	5 M	10 min
States	0	1																	
q_0	q_1	q_3																	
q_1	q_0	q_2																	
q_2	q_3	q_0																	
q_3	q_2	q_0																	
7.	<p>The given alphabet is $\Sigma = \{a, b\}$</p> <p>Transition Diagram:</p>  <p>Transition table — 1 M</p> <p>$\{\emptyset, \Sigma, \delta, q_0, F\}$ values — 1 M</p> <p>Valid String Derivation — 1 M.</p>	2 M	5 M	10 min															

Part C

(2 x 10 = 20 Marks)

Q No	Solution	Scheme of Marking	Max. Time required for each Question																											
8.	<p>The DFA obtained for the given NFA is</p> <table border="1"> <thead> <tr> <th>States</th> <th>0</th> <th>1</th> </tr> </thead> <tbody> <tr> <td>$\rightarrow \{P, P\}$</td> <td>$\{P, q_3\}$</td> <td>$\{P\}$</td> </tr> <tr> <td>$\{P, q_3\}$</td> <td>$\{P, q, s\}$</td> <td>$\{P, s\}$</td> </tr> <tr> <td>$\{P, q, s\}$</td> <td>$\{P, q, s, t\}$</td> <td>$\{P, t\}$</td> </tr> <tr> <td>$\{P, s\}$</td> <td>$\{P, q, s\}$</td> <td>$\{P\}$</td> </tr> <tr> <td>$*\{P, q, r, s\}$</td> <td>$\{P, q, r, s\}$</td> <td>$\{P, r, s\}$</td> </tr> <tr> <td>$*\{P, q, r\}$</td> <td>$\{P, q, r, t\}$</td> <td>$\{P, r, t\}$</td> </tr> <tr> <td>$*\{P, q, s, t\}$</td> <td>$\{P, q, s, t\}$</td> <td>$\{P, t\}$</td> </tr> <tr> <td>$*\{P, s, t\}$</td> <td>$\{P, q, r, s\}$</td> <td>$\{P, t\}$</td> </tr> </tbody> </table> <p>Transition Table</p> <p>Transition Diagram — 4 M</p> <p>Steps of conversion — 2 M</p>	States	0	1	$\rightarrow \{P, P\}$	$\{P, q_3\}$	$\{P\}$	$\{P, q_3\}$	$\{P, q, s\}$	$\{P, s\}$	$\{P, q, s\}$	$\{P, q, s, t\}$	$\{P, t\}$	$\{P, s\}$	$\{P, q, s\}$	$\{P\}$	$*\{P, q, r, s\}$	$\{P, q, r, s\}$	$\{P, r, s\}$	$*\{P, q, r\}$	$\{P, q, r, t\}$	$\{P, r, t\}$	$*\{P, q, s, t\}$	$\{P, q, s, t\}$	$\{P, t\}$	$*\{P, s, t\}$	$\{P, q, r, s\}$	$\{P, t\}$	4 M 10 M	15 min
States	0	1																												
$\rightarrow \{P, P\}$	$\{P, q_3\}$	$\{P\}$																												
$\{P, q_3\}$	$\{P, q, s\}$	$\{P, s\}$																												
$\{P, q, s\}$	$\{P, q, s, t\}$	$\{P, t\}$																												
$\{P, s\}$	$\{P, q, s\}$	$\{P\}$																												
$*\{P, q, r, s\}$	$\{P, q, r, s\}$	$\{P, r, s\}$																												
$*\{P, q, r\}$	$\{P, q, r, t\}$	$\{P, r, t\}$																												
$*\{P, q, s, t\}$	$\{P, q, s, t\}$	$\{P, t\}$																												
$*\{P, s, t\}$	$\{P, q, r, s\}$	$\{P, t\}$																												
9	<p>DFA Transition table — 4 M</p> <p>DFA Transition Diagram — 4 M</p> <p>Steps of conversion — 2 M.</p> <p>i. Eclose($\{\emptyset\}$) = $\{\emptyset\} = A$</p> <table border="1"> <thead> <tr> <th>Status</th> <th>a</th> <th>b</th> </tr> </thead> <tbody> <tr> <td>$\rightarrow \{A\}$</td> <td>$\{1\} = B$</td> <td>q_t</td> </tr> <tr> <td>$B = \{1\}$</td> <td>q_t</td> <td>$\{2, 3, 4, 6, 9\} = C$</td> </tr> <tr> <td>$C = \{2, 3, 4, 6, 9\}$</td> <td>$\{5, 8, 9, 3, 4, 6\} = D$</td> <td>$\{7, 8, 3, 4, 6, 9\} = E$</td> </tr> <tr> <td>$D = \{3, 4, 5, 8, 9\}$</td> <td>D</td> <td>E</td> </tr> <tr> <td>$E = \{3, 4, 6, 7, 8, 9\}$</td> <td>D</td> <td>E</td> </tr> </tbody> </table>	Status	a	b	$\rightarrow \{A\}$	$\{1\} = B$	q_t	$B = \{1\}$	q_t	$\{2, 3, 4, 6, 9\} = C$	$C = \{2, 3, 4, 6, 9\}$	$\{5, 8, 9, 3, 4, 6\} = D$	$\{7, 8, 3, 4, 6, 9\} = E$	$D = \{3, 4, 5, 8, 9\}$	D	E	$E = \{3, 4, 6, 7, 8, 9\}$	D	E	10	15 min									
Status	a	b																												
$\rightarrow \{A\}$	$\{1\} = B$	q_t																												
$B = \{1\}$	q_t	$\{2, 3, 4, 6, 9\} = C$																												
$C = \{2, 3, 4, 6, 9\}$	$\{5, 8, 9, 3, 4, 6\} = D$	$\{7, 8, 3, 4, 6, 9\} = E$																												
$D = \{3, 4, 5, 8, 9\}$	D	E																												
$E = \{3, 4, 6, 7, 8, 9\}$	D	E																												

Roll No.									
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**PRESIDENCY UNIVERSITY
BENGALURU**
SCHOOL OF ENGINEERING

TEST –2

Sem & AY: Odd Sem 2019-20

Date: 19.11.2019

Course Code: CSE 208

Time: 11:00 AM to 12:00 PM

Course Name: THEORY OF COMPUTATION

Max Marks: 40

Program & Sem: B.Tech. (CSE) & III Sem

Weightage: 20%

Instructions:

- (i) Read Questions carefully before answer the questions.
- (ii) Question Paper Contains three Parts.

Part A [Memory Recall Questions]

Answer all the Questions. Each Question carries two marks. (5Qx2M=10M)

1. Define CFG. (C.O.NO.3)[Knowledge]
2. Define Regular Expressions. (C.O.NO.3)[Knowledge]
3. Define PDA. (C.O.NO.4)[Knowledge]
4. Obtain Regular Expression for $L = \{0^n 1^m : n \geq 3, m \text{ is odd}\}$. (C.O.NO.3)[Knowledge]
5. Obtain CFG for $L = \{w w^R : w \in \{0, 1\}^*\}$. (C.O.NO.3)[Knowledge]

Part B [Thought Provoking Questions]

Answer both the Questions. Each question carries ten marks. (2Qx10M=20 M)

6. Peter want to construct a PDA to accept the Language $L = \{w \mid w \in \{a, b\}^* \text{ and } n_a(w) = n_b(w)\}$. (C.O.NO.4)[Comprehension]
7. Show that Language $L = \{a^n b^n, n \geq 0\}$ is not regular using pumping lemma. (C.O.NO.3)[Comprehension]

Part C [Problem Solving Questions]

Answer both the Questions. Each question carries five marks. (2Qx5M=10M)

8. Convert the regular expression $r = (a+b)^* abb$ into NFA with ϵ transitions. (Use thomson construction). (C.O.NO.3)[Application]
9. Consider the following CFG, $S \rightarrow aS|aSbS|\epsilon$, Show that deviation for the string aab is ambiguous. (C.O.NO.3)[Application]



SCHOOL OF ENGINEERING

Date: 19-11-2019

Time: 11:00 AM to 12:00 PM

Max Marks: 40

Weightage: 20%

Semester: III

Course Code: CSE 208

Course Name: THEORY OF COMPUTATION

Extract of question distribution [outcome wise & level wise]

Q.NO	C.O.NO	Unit/Module Number/Unit /Module Title	Memory recall type [Marks allotted] Bloom's Levels			Thought provoking type [Marks allotted] Bloom's Levels			Problem Solving type [Marks allotted]			Total Marks
			K			C			A			
1,2,4, 5, 7,8,9	3	3	8			10			10			28
3.6	4	4	2			10						12
	Total Marks		10			20			20			40

K =Knowledge Level C = Comprehension Level, A = Application Level

Note: While setting all types of questions the general guideline is that about 60%

Of the questions must be such that even a below average students must be able to attempt, About 20% of the questions must be such that only above average students must be able to attempt and finally 20% of the questions must be such that only the bright students must be able to attempt.

Annexure- II: Format of Answer Scheme



SCHOOL OF ENGINEERING

SOLUTION

Date: 19-11-2019

Time: 11:00 AM to 12:00 PM

Max Marks: 40

Weightage: 20%

Semester: III

Course Code: CSE 208

Course Name: THEORY OF COMPUTATION

Part A

(5Q x 2M = 10 Marks)

Q No	Solution	Scheme of Marking	Max. Time required for each Question
1	CFG G= (V,T,S,P) P has form A→x	Definition: 2 M	2 Min
2	Regular Expression definition	Definition: 2 M	2 Min
3	PDA M=(Q,Σ, T, δ,q0,z,F)	Definition: 2 M	2 Min
4	R= 0000*(11)*1	Correct RE- 2 M	2 Min
5	S→0S0 S→1S1 S→ ε	Correct CFG- 2 M	2 Min

Part B

(2Q x 10M = 20 Marks)

Q N o	Solution	Scheme of Marking	Max. Time required for each Question
6	$\delta(\alpha_0, \alpha, z) = (\alpha_0, \alpha z)$ $\delta(\alpha_0, b, z) = (\alpha_0, bz)$ $\delta(\alpha_0, \alpha, a) = (\alpha_0, \alpha a)$ $\delta(\alpha_0, b, b) = (\alpha_0, bb)$ $\delta(\alpha_0, a, b) = (\alpha_0, \epsilon)$ $\delta(\alpha_0, b, a) = (\alpha_0, \epsilon)$ $\delta(\alpha_0, \epsilon, z) = (\alpha_1, z)$ <u>Transition diagram for the given PDA :-</u> <pre> graph LR start(()) --> q0((q0)) q0 -- "a/a" --> q0 q0 -- "b/b" --> q0 q0 -- "z/z" --> q1(((q1))) </pre>	Transitions- 5 M Transition Diagram- 2 M Components- 1 M String Acceptance- 1 M String Rejection- 1 M	15 Min

7	<ul style="list-style-type: none"> At first, we assume that L is regular and m is the number of states. Let $w = a^m b^m$. Thus $w = 2m \geq m$. By pumping lemma, let $w = xyz$, where $xy \leq m$. Let $x = a^p$, $y = a^q$, and $z = a^r b^m$, where $p + q + r = m$, $p \neq 0$, $q \neq 0$, $r \neq 0$. Thus $y \neq 0$. Let $k = 2$. Then $xy^2z = a^p a^{2q} a^r b^m$. Number of as = $(p + 2q + r) = (p + q + r) + q = m + q$ Hence, $xy^2z = a^{m+q} b^m$. Since $q \neq 0$, xy^2z is not of the form $a^m b^m$. Thus, xy^2z is not in L. Hence L is not regular. 	String Selection – 2M Application of Pumping Lemma- 6 M Steps – 2 Marks	15 Min
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Part C

(2Q x 5 M =10 Marks)

Q No	Solution	Scheme of Marking	Max. Time required for each Question
8		NFA - 3 M Steps – 2 M	10 Min
9	<p>Two Left most derivations</p> $S \Rightarrow aSbS \Rightarrow aaSbS \Rightarrow aabS \Rightarrow aab$ $S \Rightarrow aS \Rightarrow aaSbS \Rightarrow aabS \Rightarrow aab$ <p>Thus grammar is ambiguous.</p>	Derivations – 3 M Derivation Tree – 2M	10 Min



Roll No													
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**PRESIDENCY UNIVERSITY
BENGALURU**

SCHOOL OF ENGINEERING

END TERM FINAL EXAMINATION

Semester: Odd Semester: 2019 - 20

Date: 27 December 2019

Course Code: CSE 208

Time: 1:00 PM to 4:00 PM

Course Name: THEORY OF COMPUTATION

Max Marks: 80

Program & Sem: B.Tech (CSE) & III

Weightage: 40%

Instructions:

- (i) Read the all questions carefully and answer accordingly.

Part A [Memory Recall Questions]

Answer all the Questions.

(5QX2M=10M)

1. Define automata (C.O.No.1) [Knowledge]
2. Write down the definition for CNF (C.O.No.3) [Knowledge]
3. Define ambiguous (C.O.No.3) [Knowledge]
4. State pumping lemma (C.O.No.3) [Knowledge]
5. Draw the diagram for Turing machine (C.O.No.5) [Knowledge]
6. Fill in the blanks (10QX1M=10M)
 - a. V stands for in CFG (C.O.No.3) [Knowledge]
 - b. DFA has tuple (C.O.No.2) [Knowledge]
 - c. PDA stands for (C.O.No.4) [Knowledge]
 - d. CFG stands for..... (C.O.No.3) [Knowledge]
 - e. Write down the symbol for epsilon closure.... (C.O.No.2) [Knowledge]
 - f. Write down the symbol for epsilon plus... (C.O.No.2) [Knowledge]
 - g.TM stands for..... (C.O.No.5) [Knowledge]
 - h. NPDA stands for..... (C.O.No.4) [Knowledge]
 - i. In TM, blank symbol is (C.O.No.5) [Knowledge]
 - j. GNF stands for (C.O.No.3) [Knowledge]



SCHOOL OF ENGINEERING

END TERM FINAL EXAMINATION

Extract of question distribution [outcome wise & level wise]

Q.NO.	C.O.NO (% age of CO)	Unit/Module Number/Unit /Module Title	Memory recall type [Marks allotted] Bloom's Levels	Thought provoking type [Marks allotted] Bloom's Levels	Problem Solving type [Marks allotted]	Total Marks
PART A Q. NO1 2 3 4 5 6	CO 01	All the 5 modules	20			20
	CO 02					
	CO 03					
	CO 04					
	CO 05					
PART B Q.NO.2	CO 02	MODULE 02 Introduction to Finite Automata			10	10
PART B Q.NO.3	CO 03	MODULE 03 Regular Expressions and Context Free Grammars		10		10
PART B Q.NO.4	CO 04	MODULE 04 Push Down Automata			10	10

Format of Answer Scheme



SCHOOL OF ENGINEERING

SOLUTION

Semester: Odd Sem. 2019-20

Date: 27.12.2019

Course Code: CSE208 [CSE-208]

Time: 1 TO 4 PM

Course Name: THEORY OF COMPUTATION

Max Marks: 80

Program & Sem: CSE & III SEM

Weightage: 40%

Part A

(0Q x 0M = 0Marks)

Q No	Solution	Scheme of Marking	Max. Time required for each Question
1.a	An Automata is an abstract model of a digital computer. As such, every automaton includes some essential features. It has a mechanism for reading y_p . It will be assumed that the y_p is a string over a given alphabet, written on an y_p file, which the automaton can read but not change.	2	4 min
2.b	$N \cdot T \rightarrow NT * N \cdot T$ $N \cdot T \rightarrow a$	2	4 min

Part B

(0Q x 0M = 0 Marks)

Q No	Solution	Scheme of Marking	Max. Time required for each Question															
7.a.	<p>Given Transition - Table of DFA</p> <p>0 equivalent $\{q_0 q_1, q_2 q_3\} \{q_4\}$ → 1 mark</p> <p>1 equivalent $\{q_0 q_1 q_2\} \{q_3\} \{q_4\}$ → 2 mark</p> <p>2 equivalent $\{q_0 q_2\} \{q_1\} \{q_3\} \{q_4\}$</p> <p>3 equivalent $\{q_0 q_2\} \{q_1\} \{q_3\} \{q_4\}$</p> <p>minimized DFA</p> <p>minimized Transition table DFA</p> <table border="1"> <tr> <th></th> <th>a</th> <th>b</th> </tr> <tr> <td>$q_0 q_2$</td> <td>q_1</td> <td>$q_0 q_2$</td> </tr> <tr> <td>q_1</td> <td>q_1</td> <td>q_3</td> </tr> <tr> <td>q_3</td> <td>q_1</td> <td>q_4</td> </tr> <tr> <td>q_4</td> <td>q_1</td> <td>$q_0 q_2$</td> </tr> </table>		a	b	$q_0 q_2$	q_1	$q_0 q_2$	q_1	q_1	q_3	q_3	q_1	q_4	q_4	q_1	$q_0 q_2$	<p>→ 1 mark</p> <p>→ 2 mark</p> <p>→ 2 mark</p> <p>→ 2 mark</p> <p>→ 2.5 min</p>	
	a	b																
$q_0 q_2$	q_1	$q_0 q_2$																
q_1	q_1	q_3																
q_3	q_1	q_4																
q_4	q_1	$q_0 q_2$																
7.b.	<p>$s_0 = abababa \quad s_{02} = abababab$</p>	<p>→ 1 m</p> <p>→ 5 m</p> <p>→ 5 m</p>	<p>12.5 min</p>															

9

$$\delta(q_0, a, z) = (q_0, az)$$

$$\delta(q_0, a, a) = (q_0, aa)$$

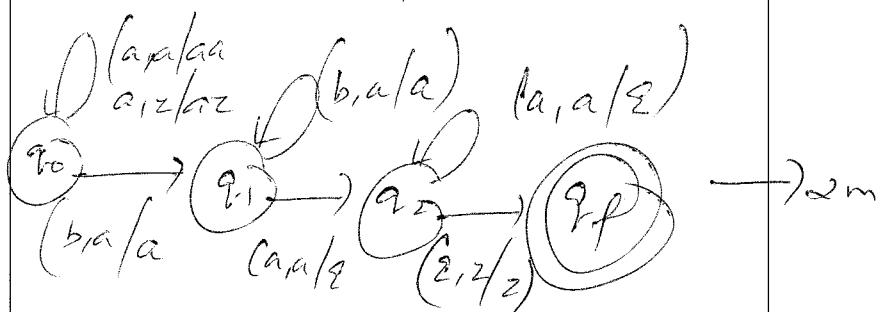
$$\delta(q_0, b, a) = (q_1, a)$$

$$\delta(q_1, b, a) = (q_1, a) \rightarrow 2m$$

$$\delta(q_1, a, a) = (q_2, z)$$

$$\delta(q_2, a, a) = (q_2, \epsilon)$$

7 tuples.



25 min

Accepting string.

Rejecting string

10m

Part C

(0Q x 0M = 0Marks)

Q No	Solution	Scheme of Marking	Max. Time required for each Question
16.a	<p>Sample Yp } Sample G/p } Tuples $n=2$ $\begin{bmatrix} aabbcc \\ xaybzc \\ xx yy zz \end{bmatrix}$</p> <p>halt</p> <p>Turing machine</p>	<p>1m</p> <p>2m</p> <p>30min</p> <p>6m</p> <p>6m</p>	